

## Homework accompanying the lecture “Basics in Applied Mathematics

### Homework 14

**Hand in:** Tuesday, 04.02.2025, after the lecture in the mailbox at the Math Institut  
(Don't forget to put your name on your homework.  
Please hand in your solutions in groups of two.)

**Exercise 1** (Heavy-Ball method for Quadratic Programming (QP); 8 points)

In this exercise, we will study the convergence rate of the Heavy-Ball method applied to some Quadratic Programming (QP) problem:

$$\min_x f(x) := \frac{1}{2}x^\top Qx - c^\top x, \quad (1)$$

with  $\mu I_n \preceq Q \preceq LI_n$  for some values  $0 < \mu < L$ .

We recall the Heavy-Ball method definition:

$$x_{k+1} = x_k - \alpha \nabla f(x_k) + \beta(x_k - x_{k-1}), \quad (2)$$

where we set  $x_{-1} = x_0$ .

We assume that  $\alpha$  and  $\beta$  are chosen such that the following condition holds:

$$\frac{(1 - \sqrt{\beta})^2}{\mu} \leq \alpha \leq \frac{(1 + \sqrt{\beta})^2}{L} \quad \text{and} \quad \beta \in (0, 1) \quad (3)$$

a) Let  $x^*$  be the solution of (1), and define the sequence  $\tilde{x}_k := x_k - x^*$ . Show that for some symmetric matrix  $A$  (that you need to find), the following holds:

$$\tilde{x}_{k+1} = A\tilde{x}_k - \beta\tilde{x}_{k-1} \quad (4)$$

b) Let  $a_1, \dots, a_n$  be the eigenvalues of  $A$ , and let  $v_1, \dots, v_n$  be a corresponding orthonormal basis of eigenvectors, i.e.:

$$Av_j = a_j v_j \quad \text{for all } i = j, \dots, n.$$

Show that for all  $j = 1, \dots, n$ , we have  $|a_j| < 2\sqrt{\beta}$ .

c) Show that:

$$\tilde{x}_k = \sum_{j=1}^n \tilde{x}_{j,k} v_j \quad (5)$$

where  $\tilde{x}_{j,k} \in \mathbb{R}$  are scalars.

In addition, express  $\tilde{x}_{j,k+1}$  as a function of  $\tilde{x}_{j,k}$  and  $\tilde{x}_{j,k-1}$ .

d) For  $j \in \{1, \dots, n\}$  and  $k \in \mathbb{N}$ , prove that:

$$z_j \tilde{x}_{j,k} - \beta \tilde{x}_{j,k-1} = (z_j)^k w_j \quad (6)$$

where  $z_j \in \mathbb{C}$  is a complex number that verifies  $z_j^2 - a_j z_j + \beta = 0$  and  $w_j \in \mathbb{C}$  is to be found.

Hint: Show that  $z_j \tilde{x}_{j,k} - \beta \tilde{x}_{j,k-1}$  is a geometric sequence.

e) Show that:

$$|\tilde{x}_{j,k}| \leq \left(\sqrt{\beta}\right)^k c_j |\tilde{x}_{0,j}| \quad (7)$$

where  $c_j$  is a constant that you need to find.

Hint 1: After ensuring that  $\text{Im}(z_j) \neq 0$ , express  $\tilde{x}_{j,k}$  in terms of  $\text{Im}(z_j \tilde{x}_{j,k} - \beta \tilde{x}_{j,k-1})$

( $\text{Im}(z)$  denotes the imaginary part of a  $z \in \mathbb{C}$ , not the image of a function!).

Hint 2: Show that  $|z_j| = \sqrt{\beta}$ .

f) Deduce that:

$$\|x_k - x^*\|^2 \leq \beta^k C \|x_0 - x^*\|^2 \quad (8)$$

for some constant  $C$  that you need to find.

g) Conclude the following:

$$f(x_k) - f(x^*) \leq \beta^k \tilde{C} (f(x_0) - f(x^*)) \quad (9)$$

for some constant  $\tilde{C}$  that you need to find.

Hint: Show that for all  $x$ :

$$\frac{\mu}{2} \|x - x^*\|^2 \leq f(x) - f(x^*) \leq \frac{L}{2} \|x - x^*\|^2$$

h) Now assume that  $\alpha$  and  $\beta$  are chosen as follows:

$$\alpha = \left(\frac{2}{\sqrt{L} + \sqrt{\mu}}\right)^2 \quad \beta = \frac{\sqrt{L} - \sqrt{\mu}}{\sqrt{L} + \sqrt{\mu}}. \quad (10)$$

Show that the assumption (3) holds for this choice of  $\alpha$  and  $\beta$ .

## Exercise 2

(A function that is difficult to optimize; 4 points)

In this exercise, we consider the following optimization problem:

$$\min_{x \in \mathbb{R}^n} f(x) := \frac{1}{2} \left( (x^{[1]} - 1)^2 + \sum_{i=1}^{n-1} (x^{[i+1]} - x^{[i]})^2 \right) \quad (11)$$

where  $x^{[i]}$  denotes the indices of  $x$  (to not mistake it with the indices of the optimization algorithm).

We consider an optimization algorithm of the following form:

$$x_{k+1} = x_k + \sum_{j=1}^k \nu_{k,j} \nabla f(x_{k-j}) \quad (12)$$

for some values of  $\nu_{k,j}$  (that might depend on the previous iterations).

a) Prove the the algorithms listed below fall into the general form (12).

- Gradient Descent
- Heavy-Ball method
- Conjugate Gradient

b) Let us assume that  $x_0 = \begin{bmatrix} 0 \\ \vdots \\ 0 \end{bmatrix}$ . Then, prove the following property for all  $k \leq n$ :

$$x_k^{[i]} = 0 \quad \text{for all } i > k. \quad (13)$$

c) Show that the problem (11) has a unique solution  $x^*$  (that needs to be found).

Prove the following inequality holds for all  $k \leq n$ :

$$\|x_k - x^*\| \geq \sqrt{1 - \frac{k}{n}} \|x_0 - x^*\| \quad (14)$$

d) To solve a convex QP, among algorithms of the form of (12), which algorithm is the fastest for finding the exact solution (in the worst-case scenario)?

### Exercise 3

(Programming exercise; 6 points)

Open the jupyter notebook, and fill in the missing parts of the code.